# Recent Advances on Generalization Bounds Part II: Combinatorial Bounds

Konstantin Vorontsov

Computing Center RAS • Moscow Institute of Physics and Technology

4th International Conference on Pattern Recognition and Machine Intelligence (PReMI'11) Moscow, Russian Federation • June 27 – July 1, 2011

## Contents

# Combinatorial framework for generalization bounds

- Probability of overfitting
- Weak (permutational) probabilistic assumptions
- OC-bound and VC-bound

# 2 Splitting and Connectivity (SC-) bounds

- SC-graph, UC-bound and SC-bound
- SC-bound is exact for some model sets of classifiers
- Proof technique: generating and inhibiting subsets

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

# Learning with binary loss

$$\mathbb{X}^{L} = \{x_{1}, \dots, x_{L}\}$$
 — a finite universe set of objects;  
 $A = \{a_{1}, \dots, a_{D}\}$  — a finite set of classifiers;  
 $I(a, x) = [$ classifier  $a$  makes an error on object  $x]$  — binary loss;

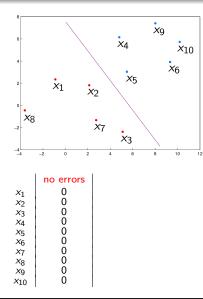
Loss matrix of size  $L \times D$ , all columns are distinct:

	$a_1$	$a_2$	<b>a</b> 3	$a_4$	$a_5$	$a_6$		$a_D$	
$x_1$	1	1	0	0	0	1		1	X — observable
	0	0	0	0	1	1		1	(training) sample
$x_\ell$	0	0	1	0	0	0	•••	0	of size $\ell$
$x_{\ell+1}$	0	0	0	1	1	1		0	$\bar{X}$ — hidden
	0	0	0	1	0	0		1	(testing) sample
XL	0	1	1	1	1	1	•••	0	od size $k = L - \ell$

n(a) — number of errors of a classifier a on the set  $\mathbb{X}^{L}$ ; n(a, X) — number of errors of a classifier a on a sample  $X \subset \mathbb{X}^{L}$ ;  $\nu(a, X) = n(a, X)/|X|$  — error rate of a on a sample  $X \subset \mathbb{X}^{L}$ ;

**Probability of overfitting** Weak (permutational) probabilistic assumptions OC-bound and VC-bound

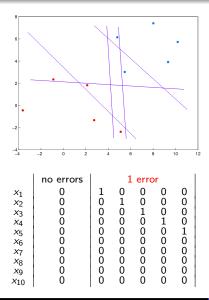
### Example. The loss matrix for a set of linear classifiers



1 vector having no errors

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

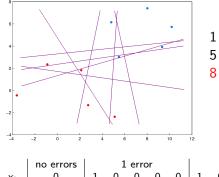
# Example. The loss matrix for a set of linear classifiers



1 vector having no errors 5 vectors having 1 error

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

### Example. The loss matrix for a set of linear classifiers



- 1 vector having no errors 5 vectors having 1 error
- 8 vectors having 2 errors

	no errors		1	erro	or					2 er	rors				
$x_1$	0	1	0	0	0	0	1	0	0	0	0	1	1	0	
$x_2$	0	0	1	0	0	0	1	1	0	0	0	0	0	0	
$x_3$	0	0	0	1	0	0	0	1	1	0	0	0	0	1	
<i>x</i> 4	0	0	0	0	1	0	0	0	1	1	0	0	0	0	
$x_5$	0	0	0	0	0	1	0	0	0	1	1	1	0	0	
$x_6$	0	0	0	0	0	0	0	0	0	0	1	0	1	0	
X7	0	0	0	0	0	0	0	0	0	0	0	0	0	1	
$x_8$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
X9	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
<i>x</i> <sub>10</sub>	0	0	0	0	0	0	0	0	0	0	0	0	0	0	

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

### Probability of overfitting

**Def.** The *learning algorithm*  $\mu \colon X \mapsto a$  takes a training sample  $X \subset \mathbb{X}^{L}$  and returns a classifier  $a \equiv \mu X \in A$ .

**Def.** Algorithm  $\mu$  overfits on a given partition  $X \sqcup \overline{X} = \mathbb{X}^{L}$  if

$$\delta(\mu, X) \equiv \nu(\mu X, \bar{X}) - \nu(\mu X, X) \ge \varepsilon.$$

### Def. Probability of overfitting

$$Q_{\varepsilon}(\mu, \mathbb{X}^{L}) = \mathsf{P}[\delta(\mu, X) \ge \varepsilon].$$

**Def.** Exact bound:  $Q_{\varepsilon} = \eta(\varepsilon)$ . **Def.** Upper bound:  $Q_{\varepsilon} \leq \eta(\varepsilon)$ .

# Weak (permutational) probabilistic assumptions

### Axiom

All partitions 
$$\mathbb{X}^{L} = \{x_{1}, \dots, x_{L}\} = X \sqcup \overline{X}$$
 are equiprobable, where  $X$  — observable training sample of size  $\ell$ ;  $\overline{X}$  — hidden testing sample of size  $k = L - \ell$ ;

**Probability** is defined as a fraction of partitions:

$$Q_{arepsilon} = \mathsf{P}ig[\delta(\mu, X) \geqslant arepsilonig] = rac{1}{C_L^\ell} \sum_{\substack{X, ar{X} \ X \sqcup ar{X} = \mathbb{X}^\ell}} ig[\delta(\mu, X) \geqslant arepsilonig].$$

**Interpretation:** Only *independence* of observations is postulated. Continuous measures, infinite sets, and limits  $|X| \rightarrow \infty$  are illegal.

Nevertheless, tight generalization bounds can be obtained!

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

# **One-classifier bound (OC-bound)**

Let 
$$A = \{a\}$$
,  $m = n(a)$ . Obviously,  $\mu X = a$  for all  $X \subset \mathbb{X}^{L}$ .

### Definition

Hypergeometric distribution function:

$$PDF: h_{L}^{\ell, m}(s) = P[n(a, X) = s] = \frac{C_{m}^{s} C_{L-m}^{\ell-m}}{C_{L}^{\ell}};$$
  

$$CDF: H_{L}^{\ell, m}(z) = P[n(a, X) \leq z] = \sum_{s=0}^{\lfloor z \rfloor} h_{L}^{\ell, m}(s)$$

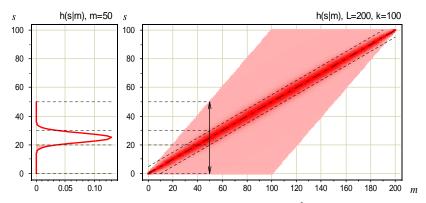
### Theorem (exact OC-bound)

For one-classifier set  $A = \{a\}$ , m = n(a), and any  $\varepsilon \in (0, 1)$ 

$$Q_{\varepsilon} = H_{L}^{\ell, m}(s_{m}(\varepsilon)), \quad s_{m}(\varepsilon) = \frac{\ell}{L}(m-\varepsilon k).$$

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

# Hypergeometric distribution, PDF $\overline{h_L^{\ell,m}(s)} = C_m^s C_{L-m}^{\ell-s} / C_L^{\ell}$



Distribution is concentrated along diagonal  $s \approx \frac{\ell}{L}m$ , thus allowing to predict both n(a) = m and  $n(a, \bar{X}) = \frac{m-s}{k}$  from n(a, X) = s. Law of Large Numbers:  $\nu(a, X) \rightarrow \nu(a)$  with  $\ell, k \rightarrow \infty$ .

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

## Vapnik-Chervonenkis bound (VC-bound), 1971

For any 
$$\mathbb{X}^L$$
,  $A$ ,  $\mu$ , and  $\varepsilon \in (0,1)$ 

$$Q_{\varepsilon} = \mathsf{P} \big[ \nu \big( \mu X, \bar{X} \big) - \nu \big( \mu X, X \big) \geqslant \varepsilon \big] \leqslant$$

**STEP 1**: *uniform bound* makes the result independent on  $\mu$ :

$$\leqslant \widetilde{Q}_{\varepsilon} = \mathsf{P}\max_{a\in A} \left[ \nu(a, \bar{X}) - \nu(a, X) \geqslant \varepsilon \right] \leqslant$$

**STEP 2**: *union bound* (wich is usually higly overestimated):

$$\leq \mathsf{P}\sum_{\mathbf{a}\in A} \left[\nu(\mathbf{a},\bar{X}) - \nu(\mathbf{a},X) \geqslant \varepsilon\right] =$$

exact one-classifier bound:

$$=\sum_{a\in A}H_L^{\ell,m}\left(s_m(\varepsilon)\right),\quad m=n(a).$$

Probability of overfitting Weak (permutational) probabilistic assumptions OC-bound and VC-bound

# OC-bound vs. VC-bound

The VC-bound [Vapnik and Chervonenkis, 1971] can be represented as a sum of OC-bounds over all classifiers  $a \in A$ :

### Theorem (OC-bound)

$$Q_{\varepsilon} = H_L^{\ell, m}(s_m(\varepsilon)), \quad m = n(a).$$

## Theorem (VC-bound)

$$Q_{\varepsilon} \leqslant \widetilde{Q}_{\varepsilon} \leqslant \sum_{a \in A} H_{L}^{\ell, m}(s_{m}(\varepsilon)), \quad m = n(a).$$

VC-bound is loose because of uniform bound and union bound, which discards the *splitting* and *similarity* properties of *A*.

# Paradigms of COLT not using union bound

- Uniform convergence bounds [Vapnik, Chervonenkis, 1968]
- Theory of learnable (PAC-learning) [Valiant, 1982]
- Data-dependent bounds [Haussler, 1992]
- Concentration inequalities [Talagrand, 1995]
- Connected function classes [Sill, 1995]
- Similar classifiers VC bounds [Bax, 1997]
- Margin based bounds [Bartlett, 1998]
- Self-bounding learning algorithms [Freund, 1998]
- Rademacher complexity [Koltchinskii, 1998]
- Adaptive microchoice bounds [Langford, Blum, 2001]
- Algorithmic stability [Bousquet, Elisseeff, 2002]
- Algorithmic luckiness [Herbrich, Williamson, 2002]
- Shell bounds [Langford, 2002]
- PAC-Bayes bounds [McAllester, 1999; Langford, 2005]
- Splitting and connectivity bounds [Vorontsov, 2010]

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

# Splitting and Connectivity graph

**Define** two binary relations on classifiers: partial order  $a \leq b$ :  $I(a, x) \leq I(b, x)$  for all  $x \in \mathbb{X}^{L}$ ; precedence  $a \prec b$ :  $a \leq b$  and Hamming distance ||b - a|| = 1.

# Definition (SC-graph)

Splitting and Connectivity (SC-) graph  $\langle A, E \rangle$ : A - a set of classifiers with distinct binary loss vectors;  $E = \{(a, b): a \prec b\}.$ 

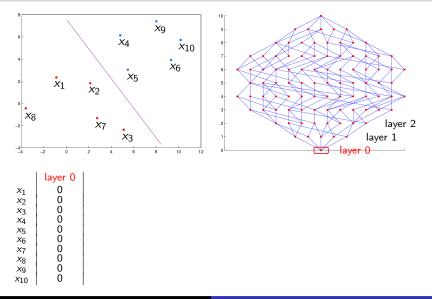
# Properties of the SC-graph:

- each edge (a, b) is labeled by an object  $x_{ab} \in \mathbb{X}^{L}$  such that  $0 = I(a, x_{ab}) < I(b, x_{ab}) = 1$ ;
- multipartite graph with layers  $A_m = \{a \in A : n(a) = m\}, m = 0, \dots, L + 1;$

SC-graph, UC-bound and SC-bound

SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

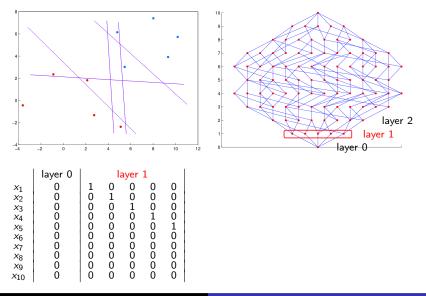
### Example. Loss matrix and SC-graph for a set of linear classifiers



SC-graph, UC-bound and SC-bound

SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

## Example. Loss matrix and SC-graph for a set of linear classifiers

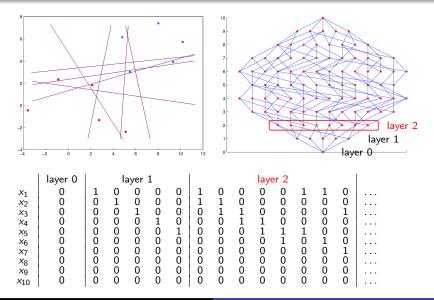


Konstantin Vorontsov www.ccas.ru/voron Recent Advances on Generalization Bounds – Part II

SC-graph, UC-bound and SC-bound

SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

## Example. Loss matrix and SC-graph for a set of linear classifiers



Konstantin Vorontsov www.ccas.ru/voron Recent Advances on Generalization Bounds – Part II

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

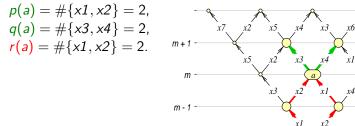
x1

### Connectivity and inferiority of a classifier

**Def.** Connectivity of a classifier 
$$a \in A$$
  
 $p(a) = \# \{ x_{ba} \in \mathbb{X}^{L} : b \prec a \}$  — low-connectivity.  
 $q(a) = \# \{ x_{ab} \in \mathbb{X}^{L} : a \prec b \}$  — up-connectivity;

# **Def.** Inferiority of a classifier $a \in A$ $r(a) = \# \{ x_{cb} \in \mathbb{X}^L : c \prec b \leq a \} \in \{ p(a), \dots, n(a) \}.$

#### Example:



SC-graph, UC-bound and SC-bound

SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

# Uniform Connectivity (UC-) bound

# Theorem (UC-bound)

For all 
$$\mathbb{X}^{L}$$
,  $\mu$ ,  $A$  and  $\varepsilon \in (0, 1)$   

$$\widetilde{Q}_{\varepsilon} \leqslant \sum_{a \in A} \left( \frac{C_{L-q-p}^{\ell-q}}{C_{L}^{\ell}} \right) H_{L-q-p}^{\ell-q, m-p}(s_{m}(\varepsilon))$$
where  $m = p(z)$ ,  $q = q(z)$ ,  $p = p(z)$ 

where m = n(a), q = q(a), p = p(a).

- UC-bound improves the VC-bound, even if  $p(a) \equiv q(a) \equiv 0$ :  $\widetilde{Q}_{\varepsilon} \leq \sum_{a \in A} H_L^{\ell, m}(s_m(\varepsilon)).$
- Provide the contribution of a ∈ A decreases exponentially by p(a)
   ⇒ connected sets are less subjected to overfitting.
- OC-bound relies on connectivity, but disregards splitting.

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

# Pessimistic Empirical Risk Minimization

# Definition (ERM)

Learning algorithm  $\mu$  is Empirical Risk Minimization if

$$\mu X \in A(X), \qquad A(X) = \operatorname{Arg\,min}_{a \in A} n(a, X);$$

A choice of a classifier a from A(X) is ambiguous. Pessimistic choice will result in modestly inflated upper bound.

## Definition (pessimistic ERM)

Learning algorithm  $\mu$  is pessimistic ERM if

$$\mu X = \arg \max_{a \in A(X)} n(a, \bar{X});$$

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers <u>Proof technique</u>: generating and inhibiting subsets

# The Splitting and Connectivity (SC-) bound

# Theorem (SC-bound)

For pessimistic ERM  $\mu$ , any  $\mathbb{X}^L$ , A and  $\varepsilon \in (0, 1)$ 

$$Q_{\varepsilon} \leq \sum_{a \in A} \left( \frac{C_{L-q-r}^{\ell-q}}{C_{L}^{\ell}} \right) H_{L-q-r}^{\ell-q, m-r}(s_{m}(\varepsilon)),$$

where m = n(a), q = q(a), r = r(a).

• If  $q(a) \equiv r(a) \equiv 0$  then SC-bound transforms to VC-bound:

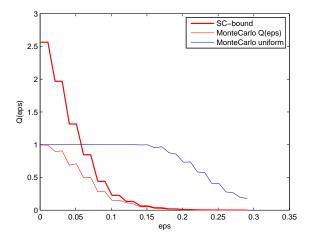
$$Q_{\varepsilon} \leq \sum_{a \in A} H_L^{\ell, m}(s_m(\varepsilon)).$$

The contribution of a ∈ A decreases exponentially by:
 q(a) ⇒ connected sets are less subjected to overfitting;
 r(a) ⇒ only lower layers contribute significantly to Q<sub>ε</sub>.

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

### Experiment on model data: SC-bound vs. Monte Carlo estimate

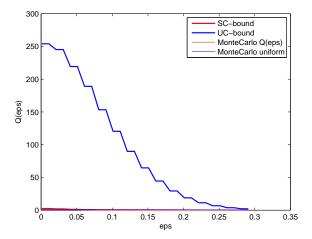
Separable two-dimensional task, L = 100, two classes.



SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

### Experiment on model data: UC-bound vs. Monte Carlo estimate

Separable two-dimensional task, L = 100, two classes.

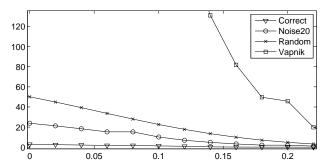


SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

### Experiment on model data: SC-bounds vs. VC-bound

Two-dimensional task, L = 100, two classes.

Correct — 0% errors; Noise20 — 20% errors; Random — 50% errors; Vapnik — data-independent VC-bound.

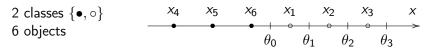


SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

### Monotone chain of classifiers

**Def.** Monotone chain of classifiers:  $a_0 \prec a_1 \prec \cdots \prec a_D$ .

**Example:** 1-dimensional threshold classifiers  $a_j(x) = [x - \theta_j]$ ;



SC-graph:

Loss matrix:

m=3 m=2 m=1 m=1 m=1 m=1 m=0 m=0 m=1

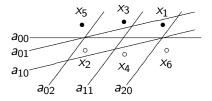
	<b>a</b> 0	$a_1$	<b>a</b> 2	<b>a</b> 3
$x_1$	0	1	1	1
$x_2$	0	0	1	1
<i>x</i> <sub>3</sub>	0	0	0	1
<i>X</i> 4	0	0	0	0
$x_5$	0	0	0	0
$x_6$	0	0	0	0

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

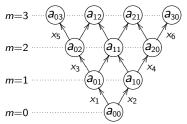
### Two-dimensional monotone lattice of classifiers

### Example:

2-dimensional linear classifiers, 2 classes {●, ○}, 6 objects



# SC-graph:



Loss matrix:

	$a_{00}$	$a_{01}$	$a_{10}$	<b>a</b> 02	$a_{11}$	<b>a</b> <sub>20</sub>	<b>a</b> 03	$a_{12}$	$a_{21}$	<b>a</b> <sub>30</sub>
$x_1$	0	1	0	1	1	0 1 0 1 0 0	1	1	1	0
<i>x</i> <sub>2</sub>	0	0	1	0	1	1	0	1	1	1
<i>X</i> 3	0	0	0	1	0	0	1	1	0	0
<i>x</i> <sub>4</sub>	0	0	0	0	0	1	0	0	1	1
$X_5$	0	0	0	0	0	0	1	0	0	0
<i>x</i> 6	0	0	0	0	0	0	0	0	0	1

### SC-bound is exact(!) for multidimensional(!) lattices of classifiers

Denote  $\mathbf{d} = (d_1, \dots, d_h)$  an *h*-dimensional index vector,  $d_j = 0, 1, \dots$ Denote  $|\mathbf{d}| = d_1 + \dots + d_h$ .

### Definition

Monotone h-dimensional lattice of classifiers of height D:

$$A = \left\{ a_{\mathbf{d}}, \ |\mathbf{d}| \leqslant D \ \middle| \begin{array}{l} \mathbf{c} < \mathbf{d} \Rightarrow a_{\mathbf{c}} < a_{\mathbf{d}} \\ n(a_{\mathbf{d}}) = m_0 + |\mathbf{d}| \end{array} \right\}$$

### Theorem (exact SC-bound)

If A is monotone h-dimensional lattice of height D,  $D \ge k$ , and  $\mu$  is pessimistic ERM then for any  $\varepsilon \in (0, 1)$ 

$$Q_{\varepsilon} = \sum_{t=0}^{k} C_{h+t-1}^{t} \frac{C_{L-h-t}^{\ell-h}}{C_{L}^{\ell}} H_{L-h-t}^{\ell-h, m_{0}} \left(s_{m_{0}+t}(\varepsilon)\right).$$

## Sets of classifiers with known SC-bound

Model sets of classifiers with known exact SC-bound:

- monotone chains and multidimensional lattices;
- unimodal chains and multidimensional lattices;
- pencils of monotone chains;
- layers and intervals of boolean cube;
- hamming balls and their lower layers;
- some sparse subsets of multidimensional lattices;
- some sparse subsets of hamming balls;

Real sets of classifiers with known tight SC-bound:

- conjunction rules (see further);
- linear classifiers (under construction now).

# Conclusions

- Combinatorial framework can give tight and sometimes exact generalization bounds.
- OC (one-classifier) bound is exact.
- UC (uniform connectivity) bound rely on *connectivity* but neglect *splitting*.
- SC (splitting and connectivity) bound is most tight and even *exact* for monotone chains and lattices of classifiers.
- SC-bound being applied to rule induction reduces testing error of classifiers by 1-2%.

**Further:** thee appendix slides about underlying combinatorial technique for SC-bounds.

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

### Generating and inhibiting subsets of objects

## Conjecture

For any  $a \in A$  generating set  $X_a \subset \mathbb{X}^L$  and inhibiting set  $X'_a \subset \mathbb{X}^L$ exist such that if classifier  $a \in A$  is a result of learning then all objects  $X_a$  lie in the training set and all objects  $X'_a$  lie in the testing set:

 $[\mathbf{v} \mathbf{v} \mathbf{v}] < [\mathbf{v} \subset \mathbf{v}] [\mathbf{v}' \subset \mathbf{v}]$ 

$$[\mu X - a] \leqslant [X_a \subseteq X] [X_a \subseteq X].$$

$$X_a$$

$$X_a$$

$$X - \text{training}$$

$$\overline{X} - \text{testing}$$

$$\mathbb{X}^L$$

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

### Bounds based on generating and inhibiting subsets

Lemma (Probability of obtaining each of classifiers)

If **Conjecture** is true then for any  $\mu$ , X,  $a \in A$ 

$$\mathsf{P}[\mu X = a] \leqslant P_a = C_{L_a}^{\ell_a} / C_L^{\ell}.$$

where  $L_a = L - |X_a| - |X'_a|$ ,  $\ell_a = \ell - |X_a|$ .

### Theorem (Probability of overfitting)

If **Conjecture** is true then for any  $\mathbb{X}^L$ ,  $\mu$ , A and  $\varepsilon \in (0, 1)$ 

$$Q_{\varepsilon} \leqslant \sum_{a \in A} P_{a} H_{L_{a}}^{\ell_{a}, m_{a}} \left( s_{a}(\varepsilon) \right),$$

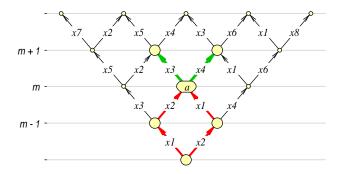
where  $m_a = n(a, \mathbb{X}^L) - n(a, X_a) - n(a, X'_a),$  $s_a(\varepsilon) = \frac{\ell}{L}(n(a, \mathbb{X}^L) - \varepsilon k) - n(a, X_a).$ 

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

Correspondence between SC-graph and generating/inhibiting subsets

Upper connectivity of a classifier  $a \in A$  $q(a) = |X_a|, X_a = \{x_{ab} \in \mathbb{X}^L : a \prec b\}$  — generating subset.

Inferiority of a classifier  $a \in A$  $r(a) = |X'_a|, X'_a = \{x_{cb} \in \mathbb{X}^L : c \prec b \leq a\}$  — inhibiting subset.



**Questions?** 

SC-graph, UC-bound and SC-bound SC-bound is exact for some model sets of classifiers Proof technique: generating and inhibiting subsets

Konstantin Vorontsov vokov@forecsys.ru http://www.ccas.ru/voron

www.MachineLearning.ru/wiki (in Russian):

- Участник:Vokov
- Слабая вероятностная аксиоматика
- Расслоение и сходство алгоритмов (виртуальный семинар)